

2010

# Good codes from generalised algebraic geometry codes

Jibril, M

<http://hdl.handle.net/10026.1/1476>

---

10.1109/ISIT.2010.5513687

Information Theory Proceedings (ISIT), 2010 IEEE International Symposium on IEEE

---

*All content in PEARL is protected by copyright law. Author manuscripts are made available in accordance with publisher policies. Please cite only the published version using the details provided on the item record or document. In the absence of an open licence (e.g. Creative Commons), permissions for further reuse of content should be sought from the publisher or author.*

# Good Codes From Generalised Algebraic Geometry Codes

Mubarak Jibril\*, Martin Tomlinson\*, Mohammed Zaki Ahmed\* and Cen Tjhai\*

\*School of Computing and Mathematics

Faculty of Technology

University of Plymouth

United Kingdom

Email: mubarak.jibril@plymouth.ac.uk

**Abstract**—Algebraic geometry codes or Goppa codes are defined with places of degree one. In constructing generalised algebraic geometry codes places of higher degree are used. In this paper we present 41 new codes over  $\mathbb{F}_{16}$  which improve on the best known codes of the same length and rate. The construction method uses places of small degree with a technique originally published over 10 years ago for the construction of generalised algebraic geometry codes.

## I. INTRODUCTION

In coding theory, it is desirable to obtain an error correcting code with the maximum possible minimum distance  $d$ , given a code length  $n$  and code dimension  $k$ . Algebraic geometry (AG) codes have good properties and some families of these codes have been shown to be asymptotically superior as they exceed the well-known Gilbert Vashamov bound [1] when the defining finite field  $\mathbb{F}_q$  has size  $q \geq 49$  with  $q$  always a square. A closer look at tables of best known codes in [2] and [3] shows that algebraic geometry codes feature as the best known linear codes for an appreciable range of code lengths for different field sizes  $q$ . Algebraic geometry codes are codes derived from curves and were first discovered by Goppa [4] in 1981. Goppa's description uses rational places of the curve to define these codes. Rational places are called places of degree one. A generalised construction of algebraic geometry codes was presented by Xing *et al* in [5] [6] and Ozbudak *et al* in [7]. An extension of the method which utilises places of higher degrees as well as a concatenation concept was introduced in [8]. This method was shown in [9] [10] [11] to be effective in constructing codes that are better than the best known codes and many codes were presented for finite fields up to  $\mathbb{F}_9$ . In this paper we present several new codes over finite field  $\mathbb{F}_{16}$ . These codes represent improvements on minimum distance compared to some previously best known codes. We first give a description of the codes and an exposition of the construction in the Section II. Finally we present our results in Section III.

## II. CONSTRUCTION

A two dimensional affine space  $\mathbb{A}^2(\mathbb{F}_q)$  is given by the set of points  $\{(\alpha, \beta) : \alpha, \beta \in \mathbb{F}_q\}$  while its projective closure  $\mathbb{P}^2(\mathbb{F}_q)$  is given by the set of equivalence points  $\{(\alpha : \beta :$

$1)\} \cup \{(\alpha : 1 : 0)\} \cup \{(1 : 0 : 0)\} : \alpha, \beta \in \mathbb{F}_q\}$ . Given a homogeneous polynomial  $F(x, y, z)$ , a curve  $\mathcal{X}$  defined in  $\mathbb{P}^2(\mathbb{F}_q)$  is a set of distinct points  $\{P \in \mathbb{P}^2(\mathbb{F}_q) : F(P) = 0\}$ . We are only interested in the case where  $\mathcal{X}$  is irreducible and is non-singular in order to obtain AG codes. Let  $\mathbb{F}_{q^\ell}$  be an extension of the field  $\mathbb{F}_q$ , the Frobenius automorphism is given as

$$\begin{aligned} \phi_{q,\ell} : \mathbb{F}_{q^\ell} &\rightarrow \mathbb{F}_{q^\ell} \\ \phi_{q,\ell}(\beta) &= \beta^q \qquad \beta \in \mathbb{F}_{q^\ell} \end{aligned}$$

and its action on a projective point  $(x : y : z)$  in  $\mathbb{F}_{q^\ell}$  is

$$\phi_{q,\ell}((x : y : z)) = (x^q : y^q : z^q).$$

A place of degree  $\ell$  [12] is a set of  $\ell$  points of a curve defined in the extension field  $\mathbb{F}_{q^\ell}$  denoted by  $\{P_0, P_1, \dots, P_{\ell-1}\}$  where each  $P_i = \phi_{q,\ell}^i(P_0)$ . Places of degree one are called rational places. An example of a place of degree two is a pair of points  $\{P_0, P_1\}$  such that  $P_0 = (x, y)$  has coordinates in  $\mathbb{F}_{q^2}$  and  $P_1 = \phi_{q,2}(P_0) = (x^q, y^q)$ .

We will now describe two maps that are useful in the Xing *et al* construction of generalised AG codes. We observe that  $\mathbb{F}_q$  is a subfield of  $\mathbb{F}_{q^\ell}$  for all  $\ell \geq 2$ . It is then possible to map  $\mathbb{F}_{q^\ell}$  to an  $\ell$ -dimensional vector space with elements from  $\mathbb{F}_q$  using a suitable basis. We define the mapping,

$$\begin{aligned} \pi_\ell : \mathbb{F}_{q^\ell} &\rightarrow \mathbb{F}_q^\ell \\ \pi_\ell(\beta_j) &= [c_1^j \ c_2^j \ \dots \ c_\ell^j] \quad \beta_j \in \mathbb{F}_{q^\ell}, \ c_i^j \in \mathbb{F}_q. \end{aligned}$$

Suppose  $[\gamma_1 \gamma_2 \dots \gamma_\ell]$  forms a suitable basis of the vector space  $\mathbb{F}_q^\ell$ , then  $\beta_j = c_1^j \gamma_1 + c_2^j \gamma_2 + \dots + c_\ell^j \gamma_\ell$ . Finally we use  $\sigma_{\ell,n}$  to represent an encoding map from an  $\ell$ -dimensional message space in  $\mathbb{F}_q$  to an  $n$ -dimensional code space,

$$\sigma_{\ell,n} : \mathbb{F}_q^\ell \rightarrow \mathbb{F}_q^n$$

with  $\ell \leq n$ .

We now give a description of generalised AG codes as presented in [8] [10] [11]. Let  $F = F(x, y, z)$  be a homogeneous polynomial defined in  $\mathbb{F}_q$ . Let  $g$  be the genus of the curve  $\mathcal{X}/\mathbb{F}_q$  corresponding to the polynomial  $F$ . Also let  $P_1, P_2, \dots, P_r$  be  $r$  distinct places of  $\mathcal{X}/\mathbb{F}_q$  and  $k_i = \deg(P_i)$

(deg is degree of).  $W$  is a divisor of the curve  $\mathcal{X}/\mathbb{F}_q$  such that  $W = P_1 + P_2 + \dots + P_r$  and  $G$  a divisor so that  $\text{supp}(W) \cap \text{supp}(G) = \emptyset$ . More specifically  $G = m(Q - R)$  where  $\text{deg}(Q) = \text{deg}(R) + 1$ . Associated with the divisor  $G$  is a Riemann-Roch space  $\mathcal{L}(G)$  with  $m = \text{deg}(G)$  an integer,  $m \geq 0$ . From the Riemann-Roch theorem we know that the dimension of  $\mathcal{L}(G)$  is given by  $l(G)$  and

$$l(G) \geq m - g + 1$$

with equality when  $m \geq 2g - 1$ . Also associated with each  $P_i$  is a  $q$ -ary code  $C_i$  with parameters  $[n_i, k_i = \text{deg}(P_i), d_i]_q$  with the restriction that  $d_i \leq k_i$ . We denote  $\{f_1, f_2, \dots, f_k : f_i \in \mathcal{L}(G)\}$  as a set of  $k$  linearly independent elements of  $\mathcal{L}(G)$  that form a basis. We can create a generator matrix for a generalised AG code as such,

$$M = \begin{bmatrix} \sigma_{k_1, n_1}(\pi_{k_1}(f_1(P_1))) & \dots & \sigma_{k_r, n_r}(\pi_{k_r}(f_1(P_r))) \\ \sigma_{k_1, n_1}(\pi_{k_1}(f_2(P_1))) & \dots & \sigma_{k_r, n_r}(\pi_{k_r}(f_2(P_r))) \\ \vdots & & \vdots \\ \sigma_{k_1, n_1}(\pi_{k_1}(f_k(P_1))) & \dots & \sigma_{k_r, n_r}(\pi_{k_r}(f_k(P_r))) \end{bmatrix}$$

where  $f_i(P_i)$  is an evaluation of a polynomial and basis element  $f_i$  at a point  $P_i$ ,  $\pi_{k_i}$  is a mapping from  $\mathbb{F}_q^{k_i}$  to  $\mathbb{F}_q$  and  $\sigma_{k_i, n_i}$  is the encoding of a message vector in  $\mathbb{F}_q^{k_i}$  to a code vector in  $\mathbb{F}_q^{n_i}$ . It is desirable to choose the maximum possible minimum distance for all codes  $C_i$  so that  $d_i = k_i$ . The same code is used in the map  $\sigma_{k_i, n_i}$  for all points of the same degree  $k_i$  i.e. the code  $C_j$  has parameters  $[n_j, j, d_j]_q$  for a place of degree  $j$ . Let  $A_j$  be an integer denoting the number of places of degree  $j$  and  $B_j$  be an integer such that  $0 \leq B_j \leq A_j$ . If  $t$  is the maximum degree of any place  $P_i$  we choose to use in the construction, then the generalised AG code is represented as a  $C(k; t; B_1, B_2, \dots, B_t; d_1, d_2, \dots, d_t)$ . Let  $[n, k, d]_q$  represent a linear code in  $\mathbb{F}_q$  with length  $n$ , dimension  $k$  and minimum distance  $d$ , then a generalised AG code is given by the parameters [8],

$$k = l(G) \geq m - g + 1$$

$$n = \sum_{i=1}^r n_i = \sum_{j=1}^t B_j n_j$$

$$d \geq \sum_{i=1}^r d_i - g - k + 1 = \sum_{j=1}^t B_j d_j - g - k + 1.$$

### III. RESULTS

We use two polynomials and their associated curves to obtain codes in  $\mathbb{F}_{16}$  better than the best known codes in [3]. The two polynomials are given in Table I while Table II gives a summary of the properties of their associated curves (with  $t = 4$ ). The number of places of degree  $j$ ,  $A_j$ , is determined by computer algebra system MAGMA [13]. The best known linear codes from [3] over  $\mathbb{F}_{16}$  with  $j = d_j$  for  $1 \leq j \leq 4$  are

$$[1, 1, 1]_{16} \quad [3, 2, 2]_{16} \quad [5, 3, 3]_{16} \quad [7, 4, 4]_{16}$$

which correspond to  $C_1$ ,  $C_2$ ,  $C_3$  and  $C_4$  respectively. Since  $t = 4$  for all the codes in this paper and

$$[d_1, d_2, d_3, d_4] = [1, 2, 3, 4]$$

TABLE I  
POLYNOMIALS IN  $\mathbb{F}_{16}$

$F_1 = x^5 z^{10} + x^3 z^{12} + x z^{14} + y^{15}$
$F_2 = x^5 + y^4 z + y z^4$

TABLE II  
PROPERTIES OF  $\mathcal{X}_i/\mathbb{F}_{16}$

$F(x, y, z)$	Genus	$A_1$	$A_2$	$A_3$	$A_4$	Reference
$\mathcal{X}_1$	12	83	60	1320	16140	[15]
$\mathcal{X}_2$	6	65	0	1600	15600	

TABLE III  
BEST CONSTRUCTIBLE CODES FROM  $\mathcal{X}_1$

Codes	$k$ Range	Description	#
$[83, k, d \geq 72 - k]_{16}$	$8 \leq k \leq 52$	$C(k; [83, 0, 0, 0])$	45
$[89, k, d \geq 76 - k]_{16}$	$9 \leq k \leq 54$	$C(k; [83, 2, 0, 0])$	46
$[94, k, d \geq 79 - k]_{16}$	$10 \leq k \leq 57$	$C(k; [83, 2, 1, 0])$	48
$[92, k, d \geq 78 - k]_{16}$	$9 \leq k \leq 57$	$C(k; [83, 3, 0, 0])$	49
$[98, k, d \geq 82 - k]_{16}$	$11 \leq k \leq 59$	$C(k; [83, 5, 0, 0])$	49

TABLE IV  
BEST CONSTRUCTIBLE CODES FROM  $\mathcal{X}_2$

Codes	$k$ Range	Description	#
$[72, k, d \geq 64 - k]_{16}$	$11 \leq k \leq 50$	$C(k; [65, 0, 0, 1])$	40
$[79, k, d \geq 68 - k]_{16}$	$11 \leq k \leq 48$	$C(k; [65, 0, 0, 2])$	38
$[77, k, d \geq 67 - k]_{16}$	$10 \leq k \leq 51$	$C(k; [65, 0, 1, 1])$	42
$[75, k, d \geq 66 - k]_{16}$	$9 \leq k \leq 51$	$C(k; [65, 0, 2, 0])$	43

TABLE V  
NEW CODES FROM  $\mathcal{X}_1$

Codes	$k$ Range	Description	#
$[70, k, d \geq 63 - k]_{16}$	$10 \leq k \leq 50$	$C(k; [65, 0, 1, 0])$	41

we shorten the representation

$$C(k; t; B_1, B_2, \dots, B_t; d_1, d_2, \dots, d_t) \equiv C(k; B_1, B_2, \dots, B_t).$$

Tables III-IV give codes obtained from the two curves associated with the two polynomials  $F_i$  for  $1 \leq i \leq 2$  that improve on the best constructible codes in the tables in [3]. Table V gives new codes that improve on both constructible and non-constructible codes in [3]. It is also worth noting that codes of the form  $C(k; N, 0, 0, 0)$  are simply Goppa codes (defined with only rational points). The symbol # in the Tables III-IV denotes the number of new codes from each generalised AG code  $C(k; B_1, B_2, \dots, B_t)$ . The tables in [14] contain curves known to have the most number of rational points for a given genus. Over  $\mathbb{F}_{16}$  the curve with the highest number of points with genus  $g = 12$  from [14] has 88 rational points, was constructed using class field theory and is not defined by an explicit polynomial. On the other hand the curve  $\mathcal{X}_1/\mathbb{F}_{16}$  obtained by Kummer covering of the projective line in [15] has  $A_1 = 83$  rational points and genus  $g = 12$  and is explicitly presented. Codes from this curve represent the best constructive codes in  $\mathbb{F}_{16}$  with code length 83. The curve  $\mathcal{X}_2/\mathbb{F}_{16}$  is defined by the well-known Hermitian polynomial.

Table VI gives the new codes obtained from  $\mathcal{X}_2$ . The codes have length  $N$ , dimension  $K$  and minimum distance  $D$ .  $D_m$

TABLE VI  
NEW CODES IN  $\mathbb{F}_{16}$

$N$	$K$	$D$	$D_m$
70	10	53	52
70	11	52	51
70	12	51	50
70	13	50	49
70	14	49	48
70	15	48	47
70	16	47	46
70	17	46	45
70	18	45	44
70	19	44	43
70	20	43	42
70	21	42	41
70	22	41	40
70	23	40	39
70	24	39	38
70	25	38	37
70	26	37	36
70	27	36	35
70	28	35	34
70	29	34	33
70	30	33	32
70	31	32	31
70	32	31	30
70	33	30	29
70	34	29	28
70	35	28	27
70	36	27	26
70	37	26	25
70	38	25	24
70	39	24	23
70	40	23	22
70	41	22	21
70	42	21	20
70	43	20	19
70	44	19	18
70	45	18	17
70	46	17	16
70	47	16	15
70	48	15	14
70	49	14	13
70	50	13	12

is the lower bound on the minimum distance of codes from [3] with the same length and dimension as the constructed generalised AG codes.

#### IV. CONCLUSION

We have presented 41 new codes and 400 improvements on constructible codes in  $\mathbb{F}_{16}$  over the codes in [3]. The construction method yields many good codes as shown here, as long as curves with many places of small degree are used, however traditional search for good codes has focused primarily on finding rational curves with many points. In order to obtain more codes from generalised AG codes, curves with small genera and many places of small degree need to be found.

#### REFERENCES

- [1] M. Tsfasman, S. Vladut, and T. Zink, "On goppa codes which are better than the varshamov-gilbert bound," *Math. Nacr.*, vol. 109, pp. 21–28, 1982.
- [2] M. Grassl, "Bounds on the minimum distance of linear codes and quantum codes," Online available at <http://www.codetables.de>, 2007, accessed on 2009-12-21.
- [3] W. Schimd and R. Shurer, "Mint: A database for optimal net parameters," Online available at <http://mint.sbg.ac.at>, 2004, accessed on 2009-12-21.
- [4] V. Goppa, *Geometry and Codes*. Dordrecht: Kluwer Academic Publishers, 1988.
- [5] C. Xing, H. Niederreiter, and K. Y. Lam, "Constructions of algebraic-geometry codes," *IEEE Trans. Inform. Theory*, vol. 45, no. 4, pp. 1186 – 1193, 1999.
- [6] H. Niederreiter, C. Xing, and K. Y. Lam, "A new construction of algebraic-geometry codes," *Appl. Algebra Engrg. Comm. Comput.*, vol. 9, no. 5, 1999.
- [7] F. Ozbudak and H. Stichtenoth, "Constructing codes from algebraic curves," *IEEE Trans. Inform. Theory*, vol. 45, no. 7, 2502-2505.
- [8] C. Xing, H. Niederreiter, and K. Lam, "A generalization of algebraic-geometry codes," *Information Theory, IEEE Transactions on*, vol. 45, no. 7, pp. 2498–2501, 1999.
- [9] K. H. Leung, S. Ling, and C. Xing, "New binary linear codes from algebraic curves," *Information Theory, IEEE Transactions on*, vol. 48, no. 1, pp. 285 –287, jan 2002.
- [10] C. Ding, H. Niederreiter, and C. Xing, "Some new codes from algebraic curves," *Information Theory, IEEE Transactions on*, vol. 46, no. 7, pp. 2638–2642, Nov 2000.
- [11] C. Xing and S. L. Yeo, "New linear codes and algebraic function fields over finite fields," *Information Theory, IEEE Transactions on*, vol. 53, no. 12, pp. 4822–4825, 2007.
- [12] J. L. Walker, *Codes and Curves*. Rhode Island: American Mathematical Society, 2000.
- [13] W. Bosma, C. J., and P. C., "The magma algebra system i: The user language," *J. Symbolic Comput.*, vol. 24, pp. 235–265, 1997.
- [14] G. van der Geer *et al.*, "Manypoints: A table of curves with many points," Online available at <http://www.manypoints.org>.
- [15] V. Shabat, "Curves with many points," Ph.D. dissertation, Univ. of Amsterdam, Amsterdam, 2001. [Online]. Available: <http://www.science.uva.nl/math/Research/Dissertations/Shabat2001.text.pdf>